Electronic Design

Proprietary RTOS APIs Impede Development

Open standard RTOS application programming interfaces, such as POSIX pthread, help enhance software development.

quick internet search shows a seemingly endless number of real-time operating systems (RTOSes), most based on proprietary application programming interfaces (APIs). That's certainly true for the more popular embedded RTOSes, e.g., <u>FreeRTOS</u>, <u>ThreadX</u>, and <u>Zephyr</u>. Some of these proprietary RTOS APIs provide adequate functionality.

However, each proprietary API impedes embedded development by requiring extensive developer training, constricting cross-platform code sharing, and effectively locking the application to the proprietary RTOS.

Developer training is time-consuming and expensive. In addition, proprietary API usage errors resulting in product defects are all too common. For device makers with MPU- (Linux) and MCU- (RTOS) based designs, sharing code between platforms is difficult, if not impossible.

Finally, since the application is locked into the RTOS, it's at the mercy of what processor and development tools are supported by that particular RTOS. Not having the ability to freely migrate your application to another hardware platform or development tool might be the most significant impediment of all.

But instead of focusing on the negative aspects of proprietary RTOS APIs, it's more productive to focus on solutions. Interestingly, the answer has been around since 1995. This is when the POSIX Threads standard (commonly called pthreads) was introduced (<u>IEEE 1003.1c-a995</u>). Not only is this an international standard, it's also the multithreading API in every embedded Linux distribution. Hence, POSIX pthreads is already the most popular API standard in the embedded industry.

Is the POSIX pthread API as functionally capable as proprietary RTOS APIs? Here's a brief overview of the POSIX pthread API alongside the most popular proprietary RTOS APIs.

Creating Threads in POSIX

Creating threads (also called tasks) is the most fundamental RTOS primitive. The POSIX *pthread create* API is ideally suited for embedded applications. Having only four arguments, it's much simpler than other popular RTOS APIs. The POSIX pthread API offers an optional attribute specification if additional configuration is required. Here are the thread creation APIs for POSIX pthreads, ThreadX, FreeRTOS, and Zephyr:

```
VOID (*thread_entry)(ULONG), ULONG thread_argument,
VOID * stack_memory, ULONG stack_size,
UINT priority, UINT preemption threshold,
```

```
ULONG time_slice, UINT auto_start);

FreeRTOS:

BaseType_t xTaskCreateStatic(TaskFunction_t pvTaskCode,

const char* const pcName,

const uint32_t ulStackDepth,

void* const pvParameters, UBaseType_t uxPriority,

StackType_t* const puxStackBuffer,

StaticTask_t* const pxTaskBuffer);

Zephyr:

k_tid_t k_thread_create(struct k_thread* new_thread,

k_thread_stack_t *stack,

size_t stack_size, k_thread_entry_t thread_entry,

void* p1, void* p2, void* p3, int prio,

uint32_t options, k_timeout_t delay);
```

USING MUTUAL EXCLUSION WITH POSIX

Mutual exclusion is required in embedded systems to coordinate access to shared resources. The mutual-exclusion primitives in POSIX are similar to the APIs in other RTOS. One difference in POSIX is that additional APIs offer timeouts for waiting on a mutex. Most proprietary RTOS APIs have an additional parameter that determines the timeout. Here are the typically used mutual-exclusion APIs for POSIX pthreads, ThreadX, FreeRTOS, and Zephyr:

```
IEEE POSIX prthread API for creating a mutex
 POSIX:
 int pthread mutex init (pthread mutex t* mutex handle,
                        pthread mutexattr t* mutex attributes);
 Proprietary RTOS API for creating a mutex
 ThreadX:
 UINT tx mutex create(TX MUTEX* mutex handle, CHAR* name,
                        UNIT inheritance option);
FreeRTOS:
 SemaphoreHandle t xSemaphoreCreateMutex(void);
 Zephyr:
 int k mutex init(struct k mutex* mutex);
 IEEE POSIX pthread API for locking a mutex
            int pthread mutex lock(pthread mutex t* mutex handle);
 Proprietary RTOS API for locking a mutex
 ThreadX:
 UINT tx mutex get(TX MUTEX* mutex handle, ULONG suspend option);
 FreeRTOS:
 UBaseType t xSemaphoreTake(SemaphoreHandle t xSemaphore,
                        TickType t xTicksToWait);
Zephyr:
 int k mutex lock(struct k mutex* mutex, k timeout t timeout);
 IEEE POSIX pthread API for unlocking a mutex
            int pthread mutex unlock(pthread mutex t* mutex handle);
```

```
Proprietary RTOS API for unlocking a mutex
ThreadX:
UINT tx_mutex_put(TX_MUTEX* mutex_handle);
FreeRTOS:
UBaseType_t xSemaphoreGive(SemaphoreHandle_t xSemaphore);
Zephyr:
int k_mutex_unlock(struct k_mutex* mutex);
```

MULTIPLE THREAD SYNCHRONIZATION USING POSIX

Synchronizing the execution of multiple threads is an important RTOS primitive. A classic example is the producerconsumer paradigm, whereby one thread processes information produced by another thread or interrupt handler.

Counting semaphores is often utilized in a producer-consumer fashion. The thread responsible for processing the information waits for a semaphore. When the information is ready, the producer sends the semaphore.

POSIX semaphore APIs are similar to proprietary RTOS semaphore APIs. One difference is that POSIX has additional APIs when a timeout is required to wait for a semaphore. Most proprietary RTOS APIs have an additional parameter that determines the timeout. Here are the commonly used semaphore APIs for POSIX, ThreadX, FreeRTOS, and Zephyr:

```
IEEE POSIX API to create a semaphore
            int sem init(sem t* semaphore handle, int pshared,
                                    unsigned int value);
 Proprietary RTOS APIs to create a semaphore
 ThreadX:
 UINT tx semaphore create(TX SEMAPHORE* semaphore handle,
                        CHAR* name, ULONG initial count);
 FreeRTOS:
 SemaphoreHandle t xSemaphoreCreateCounting(UBaseType t
      uxMaxCount, UBaseType t unInitialCount);
 Zephyr:
 int k sem init(struct k sem* sem, unsigned int initial count,
                        unsigned int limit);
 IEEE POSIX API to obtain a semaphore
            int sem wait(stm t* semaphore handle);
 Proprietary RTOS APIs to obtain a semaphore
 ThreadX:
 UINT tx semaphore get(TX SEMAPHORE* semaphore handle,
                                    ULONG suspend option);
 FreeRTOS:
 UBaseType t xSemaphoreTake(SemaphoreHandle t xSemaphore,
                                    TickType t xTicksToWait);
Zephyr:
 int k sem take(struct k sem* sem, k timeout t timeout);
```

Thread Communication in POSIX

Communicating information between multiple threads for processing is another important RTOS primitive. The thread responsible for processing the information can wait for a message from a queue. When the message is available, the waiting thread is given the message for processing and is resumed.

POSIX message-passing APIs are similar to proprietary RTOS APIs. There are differences, though. POSIX supports variable-sized messages and message priority. POSIX also has additional APIs when suspension timeouts are required. Most proprietary RTOS APIs have an additional parameter that determines the message waiting timeout. Here are the commonly used message queue APIs for POSIX, ThreadX, FreeRTOS, and Zephyr:

```
IEEE POSIX Message Queue Create API
          mqd t mq open(char* queue name, int operation, mode t mode,
                                    mq attr* attributes);
Proprietary RTOS Message Queue Create APIs
ThreadX:
UINT tx queue create(TX QUEUE* queue handle, CHAR* name,
                                    ULONG message size, VOID* queue memory,
                                    ULONG memory size);
FreeRTOS:
QueueHandle t xQueueCreateStatic(UBaseType t uxQueueLength,
                                    UBaseType t uxItemSize,
                                    uint8 t* pucQueueStorageBuffer,
                                    StaticQueue t* pxQueueBuffer);
                 void k msgq init(struct k msgq* msgq, char* buffer,
Zephyr:
                                    size t msg size, uint32 t max msgs);
IEEE POSIX Message Send API
           int mq send(mqd t queue handle, char* message,
                 size t message size, unsigned int message priority);
Proprietary RTOS Message Send API
ThreadX:
UINT tx queue send(TX QUEUE* queue handle,
                 VOID* message pointer, ULONG suspend option);
FreeRTOS:
```

```
BaseType T xQueueSend(QueueHandle t xQueue,
                  void* pvItemToQueue, TickType t xTicksToWait);
                  int k_msgq_put(struct k_msgq* msgq, const void* data, k
 Zephyr:
timeout t timeout);
 IEEE POSIX Message Receive API
            int mq receive(mqd t queue handle, char* message,
                  size t message size, unsigned int* message priority);
 Proprietary RTOS Message Receive APIs
 ThreadX:
 UINT tx queue receive(TX QUEUE* queue handle,
                  VOID* message pointer, ULONG suspend option);
 FreeRTOS:
 BaseType t xQueueReceive(QueueHandle t xQueue, void* pvBuffer,
                                     TickType t xTicksToWait);
 Zephyr:
 void* k msgq get(struct k msgq* msgq, void* data,
                                     k_timeout_t timeout);
```

POSIX PTHREAD API VS. PROPRIETARY RTOS APIs

It's clear that the POSIX pthread API is similar and just as capable as proprietary RTOS APIs for the most common multithreading primitives associated with thread management, mutual exclusion, synchronization, and message passing. The POSIX API is actually more straightforward in some cases—like creating a thread. Any missing functionality in POSIX can be augmented with API extensions.

Since most developers have some experience with POSIX pthreads, training and usage errors are significantly reduced or even eliminated. Sharing code with embedded Linux and/or moving to another RTOS that supports POSIX pthreads is easy. RTOS migration to the industry-standard IEEE POSIX pthread API promises to reduce time-to-market and enhance code reuse—welcome advances in our embedded industry.

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